# **MAEG4070** Engineering Optimization

# Lecture 1 Introduction: What is Optimization?

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## Course Setup

**Welcome** to the course on Engineering Optimization, with an introduction of its application in energy systems!

#### Basic administrative details:

#### • Instructor:

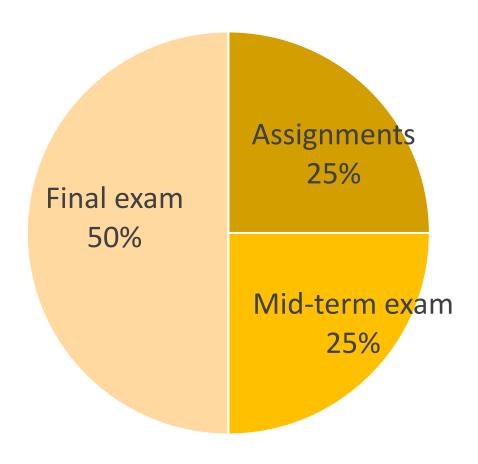
- Prof. Yue CHEN <a href="mailto:yuechen@mae.cuhk.edu.hk">yuechen@mae.cuhk.edu.hk</a> (FYB 604C)
- Office hour: 13:00-14:00, every Friday (online)
- English, Cantonese, Mandarin

## Teaching assistants:

- DENG Zhiyu <u>zydeng@mae.cuhk.edu.hk</u> (ERB113)
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# Time for tutorial? Posted on: Thursday, September 1, 2022 10:42:27 AM HKT Dear MAEG4070 students, Please select your time available for tutorial via the following link before Sep 7. https://forms.gle/zseCDZk6qdcqJFWFA The tutorial will be given by TAs, starting from week 2. Thanks! Your time available for tutorial Tuesday 13:30 - 14:15 Wednesday 10:30 - 11:15 Wednesday 12:30 - 13:15 Thursday 14:30 - 15:15 Friday 11:30 - 12:15

## **Course Grading Scheme**



- Due days for assignments will be specified.
- Assignments will carry a 50 % penalty if handed in late.
- No credit will be given to assignments which are more than three days late.
- Copying is strictly prohibited.
- No make-up exam is given unless you get approval from the Director of Registry Services for permission for absence.
- If not permitted, a zero grade in that exam will be given.

## Perquisite and reference

## Assuming work with knowledge of:

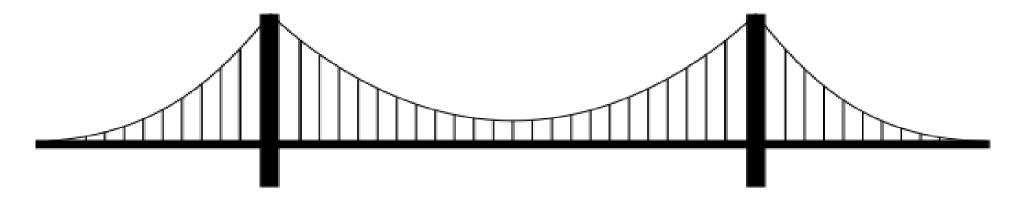
- Calculus, Linear Algebra
- Formal mathematical thinking
- Programming (Matlab, Python, ...)

If you fall short on any one of these things, don't worry!

#### Reference:

- S. Boyd, L. Vandenberghe. Convex optimization. Cambridge university press, 2004.
- R. Sioshansi, A. J. Conejo. Optimization in Engineering: Models and Algorithms. Springer, 2017.
- S. S. Rao. Engineering optimization: theory and practice. John Wiley & Sons, 2019.

## Orientation of this course



#### Basic models:

- Linear programming
- Convex optimization
- Integer programming

• .....

#### Algorithms & Solvers:

- Lagrange / Dual Methods
- Linearization approaches
- CPLEX, MOSEK, ...

### Engineering problems:

- Scheduling
- Machine Learning
- Market Analysis
- .....

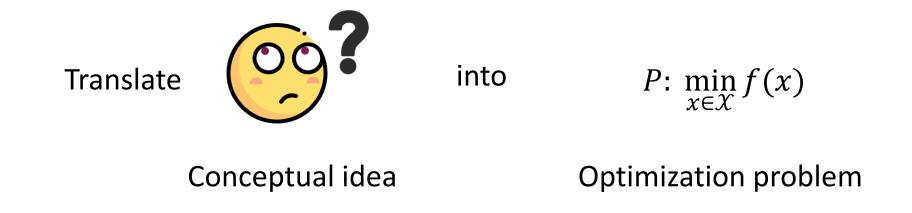
A course bridges mathematical models and engineering applications

A course combining math and programming

# What's optimization?

Different people may have different understanding about optimization

- Listing all possibilities and pick up the best?
- qualitative suggestions that lead to a better product design?
- trial and error?



A quantitative and systematic methodology to search for the best design among numerous possibilities while satisfying given constraints.

## Where can we see optimization?

As a matter of fact, any rational human behavior could be modelled as an optimization problem (intentionally or subconsciously) almost all the time.



warehouse placement

Location of warehouses

# Where can we see optimization?



Historical stock prices, economy

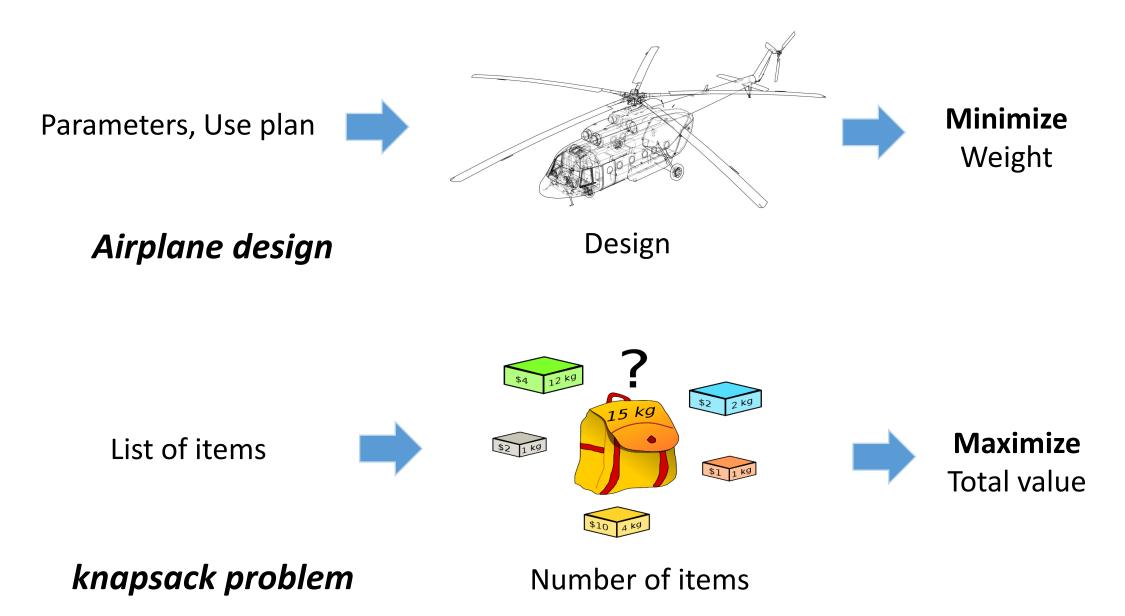


**Stock Portfolio** 



stock market

# Where can we see optimization?



# Successful Applications in Real-world

Organization	Application	Cost Saving
AT&T	Location planning of telephone selling centers	\$ 406 million
79 electric utilities	Purchasing and energy management	\$ 125 million
Delta Air Line	Aircraft scheduling	\$ 100 million
SNCF	Train scheduling	\$ 15 million

- Not all engineering problems can be quantitatively modelled as an optimization and solved with current technologies.
- With the development of more advanced mathematics and computational tools, we can model and simulate more sophisticated process, behaviors, etc.
- The application of optimization is increasingly practical, significant, and worthwhile!

Greek mathematicians considers geometrical optimization problems.

#### **Euclid**

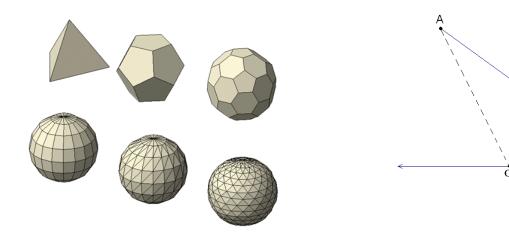
- the minimal distance between a point and a line.
- the greatest area among the rectangles with given total length of edges

#### Zenodorus

A sphere encloses the greatest volume for a given surface area

#### Heron

 light travels between two points through the path with shortest traveling time when reflected at or refracted through a boundary



## The Earliest Optimization Approach: Calculus

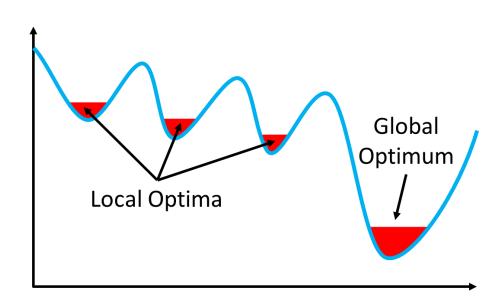
- where  $\frac{df(x)}{dx} = 0$  gives a maximum or minimum.
- Pierre De Fermat & Joseph-Louis Lagrange: calculus-based formulae for identifying optima
- Isaac Newton & Johann C.F. Gauss: iterative methods to search for an optimum.

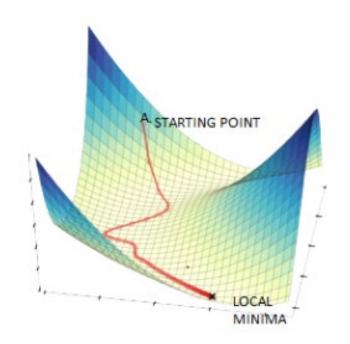
## Classic Optimization approaches

- Leonid Kantorovich: Linear Programming, 1939
- George Dantzig: the Simplex Method, 1947
- John von Neumann: the Theory of Duality, 1947
- The Karush-Kuhn-Tucker (KKT) condition, 1939 and 1951 by two separate groups
- Models: Unconstrained/Constrained optimization, single-/multi-objective optimization, etc.
- Algorithms: the steepest descent method, Newton's method, the penalty method, etc.

### First Generation Optimization: Local, Iterative, and Gradient-based

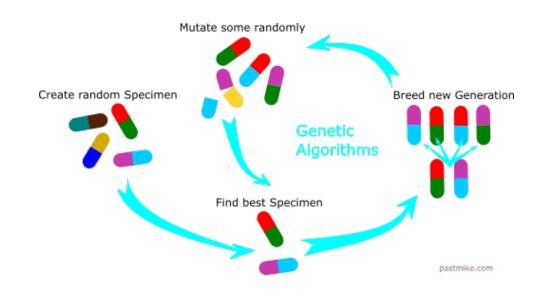
- Considers a local optimal
- Iterative methods: based on the location of the previously explored point
- Reliance on gradients or higher order derivatives
- Cons: hard to explain why optimal is optimal, limited computing power

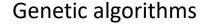


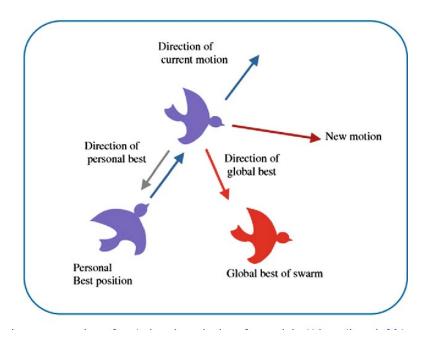


## Second Generation Optimization: The Metaheuristic Approach

- Genetic Algorithms, "survival of the fittest"; Particle Swarm Optimization
- Pros: global optimization approaches; do not require gradients; support parallel computation.
- Cons: an enormous amount of trial points; the computation time could be hours or days.



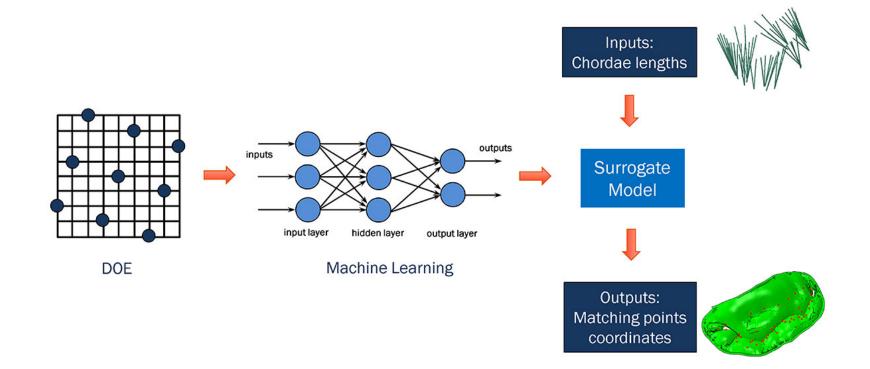




Particle swarm optimization

## Third Generation Optimization: The Al-Based Approach

- Also called surrogate-based optimization
- Pros: support parallel computation, do not need gradients, use fewer design trials than before,
   offer insight and knowledge about the design problem
- Cons: curse-of-dimensionality (Richard E. Bellman)



## Basic concept

$$P: \min_{x \in \mathcal{X}} f(x)$$

- Variable: x entities and parameters that a designer can change
- Objective: f(x) the goals that a designer wants to achieve
- Constraints:  $\mathcal{X}$  technical restrictions that a product or process must satisfy; or can be called *feasible region*; x is a feasible point if  $x \in \mathcal{X}$ .

### An optimization is either:

- Infeasible, where  $\mathcal{X} = \Phi$
- Unbounded, where  $\mathcal{X} \neq \Phi$ , but f(x) can go to infinitely negative (unbounded below, also called "optimum/minimum not attained") or infinitely positive (unbounded above, usually treated as an infeasible point) when x varies in  $\mathcal{X}$
- Feasible, where  $\mathcal{X} \neq \Phi$ , and the value f(x) is finite.

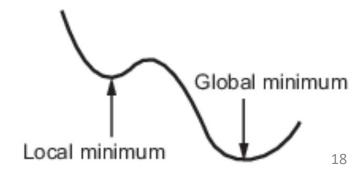
# **Basic concept**

**Global optimum.** Let f(x) be the objective function,  $\mathcal{X}$  be the feasible region, and  $x_0 \in \mathcal{X}$ . Then  $x_0$  is the global optimum if and only if  $f(x) \ge f(x_0)$ ,  $\forall x \in \mathcal{X}$ .

**Local optimum**. Let f(x) be the objective function,  $\mathcal{X}$  be the feasible region, and  $x_0 \in \mathcal{X}$ . If there is a neighborhood of  $x_0$  with radius  $\varepsilon > 0$ :

$$\mathcal{N}_{\epsilon}(x_0) = \{x \mid ||x - x_0|| < \epsilon\}$$

Such that  $\forall x \in \chi \cap N_{\varepsilon}(x_0)$ , we have  $f(x) \ge f(x_0)$ . Then  $x_0$  is a local optimum.



# Classification of optimization problems

#### **Problem:**

- Constrained v.s. nonconstrained
- Single-objective v.s. multi-objective
- Single-level v.s. multi-level
- Deterministic v.s. uncertain

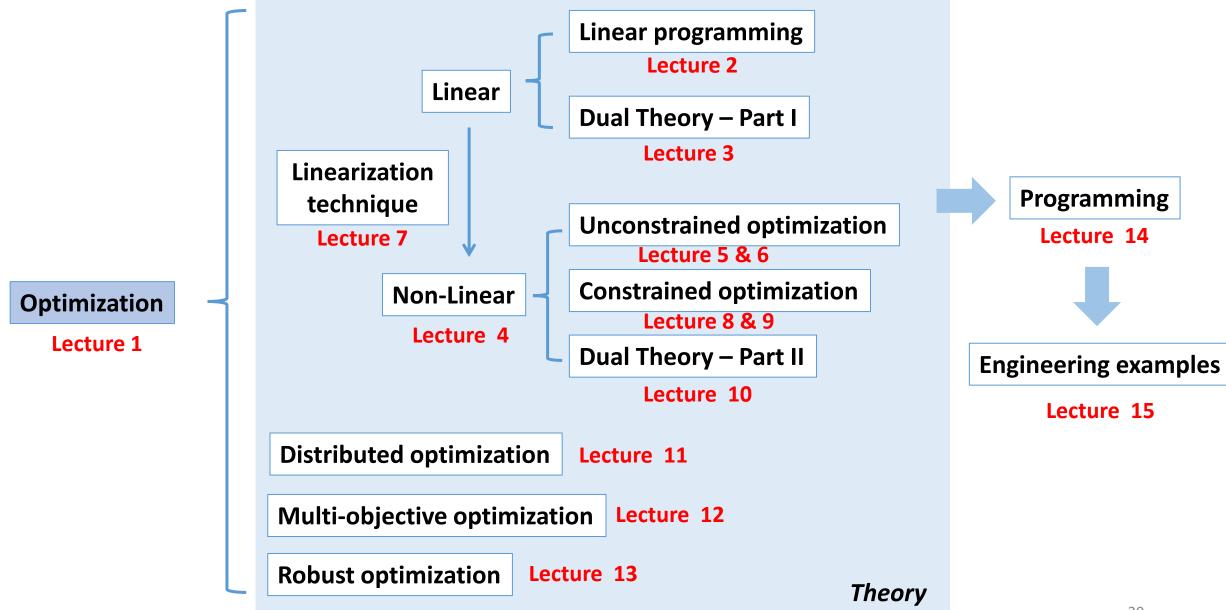
#### Response:

- Linear v.s. nonlinear
- Convex v.s. nonconvex
- Smooth v.s. nonsmooth

#### Variable:

Continuous v.s. discrete (integer/binary)

# Content of this course (tentative)



## Procedure of applying optimization in engineering

In general, applying optimization in engineering follows five steps:

Step 1: Propose optimization problems, collect relevant data and information

**Step 2:** Establish mathematical model: determine the variables, the objective function and constraints

**Step 3:** Analyze the model and select the appropriate optimization method

**Step 4:** Solve: obtain the optimal solution via programming

**Step 5:** Test and implementation of optimal solution



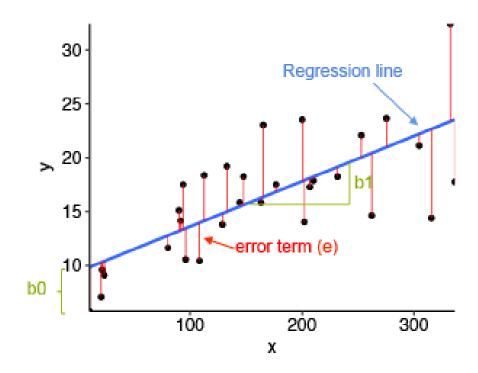
## **Example-1 Parametric Regression**

Given a data set  $\{y_i, x_i\}$ ,  $\forall i = 1, ..., N$  we assume that the relationship between the dependent variable y and the regressor x is linear.

$$y = b_0 + b_1 x + \varepsilon$$

We want to minimize the sum of square error.

$$\min_{b_0,b_1} \sum_{i=1}^{N} (y_i - b_0 - b_1 x_i)^2$$



## **Example-2 Production planning**

A company has some resources to produce three products (denoted as A, B, C). Each product consumes a different mix of resources, and there will be a profit from selling the product. The endowment of resources and its relationship with products are:

	Α	В	С	Endowment
Steel	3	4	2	600
Wood	2	1	2	400
Label	1	3	3	300
Machine	1	4	4	200
Profit	2	4	3	

Question: How to maximize the total profit?

## **Example-2 Production planning**

*Variables*: Denote  $x_1, x_2, x_3$  be the production of products A, B, C, respectively.

Objective: How to maximize the total profit?

Constraints: do not violate the resource endowment

$$\max_{x_1, x_2, x_3} 2x_1 + 4x_2 + 3x_3$$
 s.t.  $3x_1 + 4x_2 + 2x_3 \le 600$  Endowment limits 
$$2x_1 + x_2 + 3x_3 \le 400$$
 
$$x_1 + 3x_2 + 3x_3 \le 300$$
 
$$x_1 + 4x_2 + 4x_3 \le 200$$
 
$$x_1, x_2, x_3 \ge 0$$
 Production must larger than 0

## **Example-3 Transportation**

A building materials company has three cement factories  $A_1, A_2, A_3$  and four dealers  $B_1, B_2, B_3, B_4$ . Its output, sales volume and freight are shown in the table below

	$B_1$	$B_2$	$B_3$	$B_4$	Production
$A_1$	8	7	3	2	2000
$A_2$	4	7	5	1	10000
$A_3$	2	4	9	6	4000
Sales	3000	2000	4000	5000	

Question: How to minimize the freight?

## **Example-3 Transportation**

Variables: Denote  $x_{ij}$  as the quantity from factory  $A_i$  to dealer  $B_i$ 

Constraints: Production ability limits; selling requirement

Objective: Minimize the total freight

$$\begin{array}{ll} \min\limits_{x_{ij},\forall i,j} \ 8x_{11} + 7x_{12} + 3x_{13} + 2x_{14} + 4x_{21} + 7x_{22} \\ & + 5x_{23} + x_{24} + 2x_{31} + 4x_{32} + 9x_{33} + 6x_{34} \\ \mathrm{s.t.} \ x_{11} + x_{12} + x_{13} + x_{14} \leq 2000 & \mathsf{Production ability} \\ & x_{21} + x_{22} + x_{23} + x_{24} \leq 10000 & \mathsf{x}_{31} + x_{32} + x_{33} + x_{34} \leq 4000 \\ & x_{11} + x_{21} + x_{31} \geq 3000 & \mathsf{Sells requirement} \\ & x_{12} + x_{22} + x_{32} \geq 2000 & \mathsf{Sells requirement} \\ & x_{13} + x_{23} + x_{33} \geq 4000 & \mathsf{Quantity must larger than 0} \\ & x_{14} + x_{24} + x_{34} \geq 5000 & \mathsf{Quantity must larger than 0} \\ & x_{ij} \geq 0, i = 1, 2, 3; j = 1, 2, 3, 4 & \mathsf{Quantity must larger than 0} \end{array}$$

## Example-4 Location problem

Suppose there are j=1,...,n markets whose location is  $(a_j,b_j)$ . Market j needs  $q_j$  product. We plan to build m warehouse whose capacities are  $c_i$ , i=1,...,m. How to choose the location of these warehouses so that the total cost is minimized?

$$\min_{x_i, y_i, w_{ij}} \sum_{j=1}^n \sum_{i=1}^m w_{ij} \sqrt{(x_i - a_j)^2 + (y_i - b_j)^2}$$
s.t. 
$$\sum_{j=1}^n w_{ij} \le c_i, \forall i = 1, ..., m$$

$$\sum_{i=1}^m w_{ij} = q_j, \forall j = 1, ..., n$$

$$w_{ij} \ge 0, i = 1, ..., m, j = 1, ..., n$$

# Thanks!